

#k-PARTITIONS Time, Space $O^*(2^n)$

$$S \subseteq 2^U$$

множество p-pa 2^n

$$\begin{aligned} 2^n \cdot n \\ \downarrow S_0, S_1, S_2, \dots, S_n \\ N(W, j) = \text{число } S \in S, \text{ где } |S| = j \end{aligned}$$

$$S[S] = \begin{cases} 0, S \notin S \\ 1, S \in S \end{cases}$$

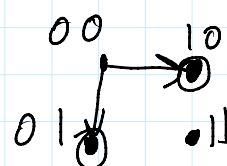
$$S'[j; S] = \begin{cases} 0, S \notin S \vee |S| \neq j \\ 1, \text{ иначе} \end{cases}$$

$$\mathcal{L}^S(S) = \sum_{X \subseteq S} S(X)$$

$$N(W, j) = \mathcal{L}^S S'_j (U|W) = \sum_{X \subseteq U|W} S(j; X)$$

Леб: \mathcal{L} -пресекование за $O^*(2^n)$

for $X \subseteq U$: [for $i \in U$:
 | for $i \in X$: for $i \notin X \subseteq U$:
 | $S[X] += S[X \setminus \{i\}]$



$$\mathcal{L}_0 S$$

||

$$S \rightarrow \mathcal{L}_1 S \rightarrow \mathcal{L}_2 S \rightarrow \mathcal{L}_3 S \rightarrow \dots \rightarrow \mathcal{L}_n S(X)$$

$$\mathcal{L}_i S(X) = \sum_{S \setminus \{i_1, i_2, \dots, i_i\} \subseteq X \subseteq S} S(X)$$

$$\mathcal{L}_i S(S) = \sum_{S \setminus \{i_1, i_2, \dots, i_i\} \subseteq X \subseteq S} S'(X) + \sum_{\text{другие}} S(X)$$

$$\sum_i S(S) = \sum_{\substack{S \subseteq \{1, 2, \dots, n\} \\ X \subseteq S \\ X \neq i}} S(X) + \sum_{\substack{S \subseteq \{1, 2, \dots, n\} \\ X \subseteq S \\ X \neq i}} S(X)$$

\downarrow

$$S \setminus \{i\} \subseteq X \subseteq S$$

\downarrow

$$\sum_{i=1}^n S(S)$$

\downarrow

$$\sum_{i=1}^n S(S \setminus \{i\}) = \sum_{i=1}^n S(S \setminus \{i\})$$

$$\sum_i S(S) = \sum_{i=1}^n S'(S), \text{ even } i \notin S$$

$$\sum_i S'(S) = \sum_{i=1}^n S(S) - \sum_{i=1}^n S(S \setminus \{i\}), \text{ even } i \in S$$

$\text{poly}(n) \cdot 2^n$

SUBSET CONVOLUTION

$$f, g: \{0, 1\}^n \rightarrow \mathbb{Q}$$

$$f * g: \{0, 1\}^n \rightarrow \mathbb{Q}$$

$$\forall S \subseteq [n] \quad (f * g)(S) = \sum_{X \subseteq S} f(X) \cdot g(S \setminus X)$$

$$f(S) = \alpha^{|S|} \quad g(S) = \beta^{|S|} \quad (f * g)([n])$$

$$\sum_k \binom{n}{k} \alpha^k \beta^{n-k}$$

$$\sum_{X \subseteq [n]} \alpha^{|X|} \beta^{|[n] \setminus X|}$$

$$f * g(S) = \sum_{X \subseteq S} f(X) \cdot g(S \setminus X)$$

$$f'(x) = (-1)^{|M|} \cdot f(x)$$

$$\mathcal{L} f(S) \cdot (-1)^{|S|} = \mu(S) \quad (-1)^{|S|} \cdot (-1)^{|X|}$$

$$\mathcal{L} f(S) \cdot \mathcal{L} g(S)$$

$$\boxed{\sum_{X, Y \subseteq S} f(X) \cdot g(Y)}$$

$X \cup Y = S$

$$\mu f(S) = \sum_{X \subseteq S} (-1)^{|S \setminus X|} \cdot f(X)$$

$$\mu f = f * (-1)^{|S|}$$

$$\mu (\mathcal{L} f) \equiv f$$

$$\mathcal{L} f = f * 1$$

$$\mu(\mathcal{Z}f) \stackrel{def}{=} f$$

$\mathcal{Z}f = + * 1$

$$\mu(\mathcal{Z}f)(S) = \sum_{X \subseteq S} (-1)^{|S| \setminus |X|} \mathcal{Z}f(X) = \sum_{X \subseteq S} (-1)^{|S| \setminus |X|} \sum_{Y \subseteq X} f(Y) =$$

$$= \sum_{Y \subseteq S} f(Y) \cdot \sum_{X \supseteq Y} (-1)^{|S| \setminus |X|} = \sum_{Y \subseteq S} f(Y) \cdot \sum_{Z \subseteq S \setminus Y} (-1)^{|S| \setminus |Y|}$$

$X = Y \cup Z$

$f(S)$

$(1 + (-1))^{|S| \setminus |Y|}$

$$\mu(\mathcal{Z}f \cdot \mathcal{Z}g)(S) = \sum_{X \subseteq S} (-1)^{|S| \setminus |X|} \mathcal{Z}f(X) \cdot \mathcal{Z}g(X) =$$

$$= \sum_{X \subseteq S} (-1)^{|S| \setminus |X|} \sum_{Y, Z \subseteq X} f(Y) \cdot g(Z)$$

$$= \sum_{Y, Z \subseteq S} f(Y) \cdot g(Z) \cdot \left(\sum_{X \supseteq Y \cup Z} (-1)^{|S| \setminus |X|} \right)$$

$$= \sum_{\substack{Y, Z \subseteq S \\ Y \cap Z = \emptyset}} f(Y) \cdot g(Z) \sum_{T \subseteq S \setminus (Y \cup Z)} (-1)^{|S \setminus (Y \cup Z)| \setminus |T|}$$

$(1 + (-1))^{|S \setminus (Y \cup Z)|}$

$$\sum_{Y \cup Z = S} f(Y) \cdot g(Z) = \mu(\mathcal{Z}f \cdot \mathcal{Z}g)(S) = f * g$$

$Y \cap Z = \emptyset$ $|Y| + |Z| = S$

$$\mu(\sum_k \mathcal{Z}f_k \cdot \mathcal{Z}g_{|S|-k})(S)$$

$$\mathcal{Z}f(k, S) = \sum_{\substack{X \subseteq S \\ |X| = k}} f(X)$$

Thm 1.6 гласио је да $f: \{-1, 1\}^n \rightarrow \{-M, \dots, M\}$

$\mu f \neq \mathcal{Z}f$ најбољи временни сложност је $\text{poly}(n) \cdot \log M \cdot 2^n$

$f \circ f$ можно вычислить за $\text{poly}(n) \cdot \log M \cdot 2^n$

Cor f^k можно вычислить за $\text{poly}(n) \cdot \log M \cdot 2^n$

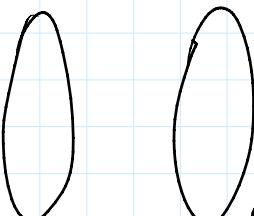
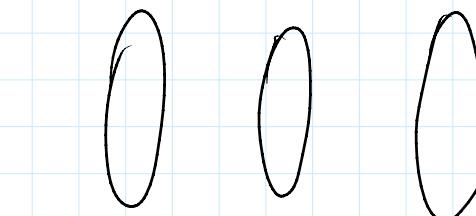
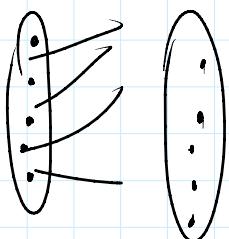
$$f : \{0,1\}^n \rightarrow \{0,1\}$$

$$f[S] = [S \in S]$$

$f * f * f * f * f(S)$ = число разбиений S на k нон- \emptyset подмножеств

f биекц.
 k раз

$$f \circ f \circ f$$



$$g(l, S) = \begin{cases} 1, & f(S) = l \\ 0, & \text{иначе} \end{cases}$$

$f(S) = \text{число разб., близкое к } \log_2 n \text{ разбиваний } S$

$$f(\emptyset) = +\infty$$

$$\min_{\substack{S_1, S_2, \dots, S_k \\ S_1 \cup S_2 \cup \dots \cup S_k = S}} f(S_1) + f(S_2) + f(S_3) + \dots + f(S_k) \rightarrow \min$$

(+, *)

\downarrow
(min, +)

$$D(k, S, l) = \begin{cases} 1, & \text{если } k \text{ ненул. и } S \text{ не разб. на } k \text{ ненул. подмнож.} \\ 0, & \text{иначе} \end{cases}$$

$$D(k, \cdot, l) = \sum_{l' \leq l} g(l', \cdot) * D(k-1, \cdot, l-l')$$

$$D(k, S, l) = \bigvee_{X \subseteq S} D(k-1, S \setminus X, l-f(X))$$

$$(f * g)(S) = \max_{X \subseteq S} [f(X) + g(S \setminus X)]$$

$$f'(S) = \beta^{f(S)}$$

$$g'(S) = \beta^{g(S)}$$

$$(f' * g') = \sum_{X \cup Y = S} \beta^{f(X) + g(Y)}$$

$$(f * g) = \log_{\beta}(f' * g')$$

$$\sum_{i=1}^n \alpha_i \beta^{f(X) + g(Y)}$$

(Reduction to $\alpha_i < \beta$)

$$\beta > \alpha_i$$

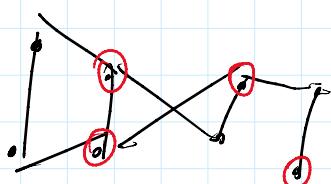
$$\beta = 2^{n+1}$$

$$f, g: \{0, 1\}^n \rightarrow \{0, \dots, M\}$$

$$f', g' \rightarrow \{1, \dots, 2^{(n+1)M}\}$$

Then \oplus (\max , $+$) B_{014} . za $2^n \cdot \text{poly}(n) \cdot M^{O(1)}$

STEINER TREE C k-так. за 2^k



Найкорочіший меж. сполучення
составлено з k дерев

$$ST(v, k, S) = \min_{\substack{\text{all trees} \\ \text{with } k \text{ components}}} \text{cost}$$

of the tree

$$ST(v, k, S) = \min_{u \in N(v)} \min_{\substack{k' \leq k, S' \subseteq S}} ST(u, k', S') + ST(v, k-k', S \setminus S')$$

$$ST(v, k): \{0, 1\}^k \rightarrow [n]$$

$\leftarrow \dots \rightarrow \dots$

$$ST(v, k) = \min_u \min_{k'} ST(u, k') \oplus ST(v, k-k')$$

$$S^I(v, k) : \{0, 1\}^S \rightarrow [n]$$

$$ST^I(v, k) = ST(v, k) + 1$$

$$ST(v, k) = \min_u \min_{k'} ST^I(u, k') \oplus ST(v, k')$$

G, k make. по разному порожденной многопар,

которому расп. в k веров

$$f(S) = 1 \Leftrightarrow S - \text{регул. МН-бо}$$

$$\underbrace{f * f * f * f}_{k} (V(G))$$

$$G[S]$$

расп. в k веров